Database Principles: Fundamentals of Design, Implementation, and Management Tenth Edition

Chapter 4 Relational Model Characteristics

Objectives

In this chapter, students will learn:

- That the relational database model offers a logical view of data
- About the relational model's basic component: relations
- That relations are logical constructs composed of rows (tuples) and columns (attributes)
- That relations are implemented as tables in a relational DBMS

Objectives (cont'd.)

- About relational database operators, the data dictionary, and the system catalog
- How data redundancy is handled in the relational database model
- Why indexing is important

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A Logical View of Data

- Relational model
 - View data logically rather than physically
- Table
 - Structural and data independence
 - Resembles a file conceptually
- Relational database model is easier to understand than hierarchical and network models

Tables and Their Characteristics

- Logical view of relational database is based on relation
 - Relation thought of as a table
- Table: two-dimensional structure composed of rows and columns
 - Persistent representation of logical relation
- Contains group of related entities (entity set)

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TABLE 4.1

Characteristics of a Relational Table

1	A table is perceived as a two-dimensional structure composed of rows and columns.
2	Each table row (tuple) represents a single entity occurrence within the entity set.
3	Each table column represents an attribute, and each column has a distinct name.
4	Each intersection of a row and column represents a single data value.
5	All values in a column must conform to the same data format.
6	Each column has a specific range of values known as the attribute domain.
7	The order of the rows and columns is immaterial to the DBMS.
8	Each table must have an attribute or combination of attributes that uniquely identifies each row.

FIGURE 4.1

STUDENT table attribute values

Table name: STUDENT

Database name: Ch04_TinyCollege

wser William C 12-Feb-1975 42 So 2.84 No BIOL 2134	
MISEL MITHENT C TEFEDERIS 44,50 2.04 NULDIOL 2134	205
nithson Anne K 15-Nov-1981 81 Jr 3.27 Yes CIS 2256	223
ewer Juliette 23-Aug-1969 36 So 2.26 Yes ACCT 2256	22
alonski Woher H 16-Sep-1976 66 Jr 3.09 No CIS 2114	22;
nith John D 30-Dec-1958 102 Sr 2.11 Yes ENGL 2231	19:
dinga Raphaol P 21-Oct-1979 114 Sr 3.15 No ACCT 2267	223
abertson Gerald T 08-Apr-1973 120 Sr 3.87 No EDU 2267	31
nith John B 30-Nov-1986 15 Fr 2.32 No ACCT 2315	23

0.00_0.0000		
STU_LNAME	= Student last name	
STU_FNAME	= Student first name	
STU_INIT	= Student middle initial	
STU_DOB	= Student date of birth	
STU_HRS	= Credit hours earned	
STU_CLASS	= Student classification	
STU_GPA	= Grade point average	
STU_TRANSFER	= Student transferred from another institution	
DEPT_CODE	= Department code	
STU_PHONE	= 4-digit campus phone extension	
PROF_NUM	= Number of the professor who is the student's advisor	
	· · · · · · · · · · · · · · · · · · ·	SOURCE: Course Technology/Cengage Learning
	STU_LNAME STU_FNAME STU_INIT STU_DOB STU_HRS STU_CLASS STU_CLASS STU_GPA STU_TRANSFER DEPT_CODE STU_PHONE	STU_LNAME= Student last nameSTU_FNAME= Student first nameSTU_INIT= Student middle initialSTU_DOB= Student date of birthSTU_HRS= Credit hours earnedSTU_CLASS= Student classificationSTU_GPA= Grade point averageSTU_TRANSFER= Student transferred from another institutionDEPT_CODE= Department codeSTU_PHONE= 4-digit campus phone extensionPROF_NUM= Number of the professor who is the student's advisor

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Keys

- Each row in a table must be uniquely identifiable
- Key: one or more attributes that determine other attributes
 - Key's role is based on determination
 - If you know the value of attribute A, you can determine the value of attribute B
 - Functional dependence
 - Attribute B is functionally dependent on A if all rows in table that agree in value for A also agree in value for B

TABLE Student Classification			
HOURS COMPLETED	CLASSIFICATION		
Less than 30	Fr		
30–59	So		
60–89	Jr		
90 or more	Sr		

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Types of Keys

- Composite key
 - Composed of more than one attribute
- Key attribute
 - Any attribute that is part of a key
- Superkey
 - Any key that uniquely identifies each row
- Candidate key
 - A superkey without unnecessary attributes

Types of Keys (cont'd.)

- Entity integrity
 - Each row (entity instance) in the table has its own unique identity
- Nulls
 - No data entry
 - Not permitted in primary key
 - Should be avoided in other attributes

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Types of Keys (cont'd.)

- Can represent:

- An unknown attribute value
- A known, but missing, attribute value
- A "not applicable" condition
- Can create problems when functions such as COUNT, AVERAGE, and SUM are used
- Can create logical problems when relational tables are linked

Types of Keys (cont'd.)

- Controlled redundancy
 - Makes the relational database work
 - Tables within the database share common attributes
 - Enables tables to be linked together
 - Multiple occurrences of values not redundant when required to make the relationship work
 - Redundancy exists only when there is unnecessary duplication of attribute values

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FIGURE An example of a simple relational database 4.2Table name: **PRODUCT** Database name: Ch04_SaleCo Primary key: PROD_CODE Foreign key: VEND_CODE PROD CODE PROD DESCRIPT PROD PRICE PROD ON HAND VEND CODE 001278-AB Claw hammer 12.95 23 232 123-21UUY Houselite chain saw, 16-in. bar 189.99 4 235 QER-34256 Sledge hammer, 16-lb. head 18.63 6 231 SRE-657UG Rat-tail file 2.9915 232 ZZX/3245Q Steel tape, 12-ft. length 679 8 235 link VEND_CODE VEND_CONTACT VEND_AREACODE VEND_PHONE Table name: VENDOR 230 Shelly K. Smithson 608 555-1234 Primary key: VEND_CODE 231 James Johnson 615 123-4536 232 Annelise Crystall 224-2134 608 Foreign key: none 233 Candice Wallace 904 342-6567 234 Arthur Jones 615 123-3324 235 Henry Ortozo 615 899-3425

SOURCE: Course Technology/Cengage Learning

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Types of Keys (cont'd.)

• Foreign key (FK)

An attribute whose values match primary key values in the related table

- Referential integrity
 - FK contains a value that refers to an existing valid tuple (row) in another relation
- Secondary key

- Key used strictly for data retrieval purposes

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TABLE 4.3

Relational Database Keys

KEY TYPE	DEFINITION
Superkey	An attribute or combination of attributes that uniquely identifies each row in a table
Candidate key	A minimal (irreducible) superkey; a superkey that does not contain a subset of attributes that is itself a superkey
Primary key	A candidate key selected to uniquely identify all other attribute values in any given row; cannot contain null entries
Foreign key	An attribute or combination of attributes in one table whose values must either match the primary key in another table or be null
Secondary key	An attribute or combination of attributes used strictly for data retrieval purposes

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Integrity Rules

- Many RDBMs enforce integrity rules automatically
- Safer to ensure that application design conforms to entity and referential integrity rules
- · Designers use flags to avoid nulls
 - Flags indicate absence of some value

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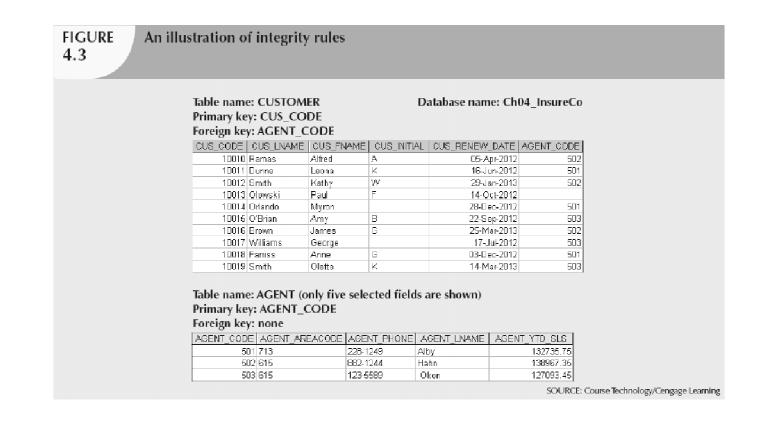
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TABLE 4.4 Integrity Rul	les
ENTITY INTEGRITY	DESCRIPTION
Requirement	All primary key entries are unique, and no part of a primary key may be null.
Purpose	Each row will have a unique identity, and foreign key values can properly reference primary key values.
Example	No invoice can have a duplicate number, nor can it be null. In short, all invoices are uniquely identified by their invoice number.
REFERENTIAL INTEGRITY	DESCRIPTION
Requirement	A foreign key may have either a null entry, as long as it is not a part of its table's primary key, or an entry that matches the primary key value in a table to which it is related. (Every non-null foreign key value <i>must</i> reference an <i>existing</i> primary key value.)
Purpose	It is possible for an attribute <i>not</i> to have a corresponding value, but it will be impossible to have an invalid entry. The enforcement of the referential integrity rule makes it impossible to delete a row in one table whose primary key has mandatory matching foreign key values in another table.
Example	A customer might not yet have an assigned sales representative (number), but it will be impossible to have an invalid sales representative (number).

TABLE A Dummy Variable Value Used as a Flag 4.5				
AGENT_CODE	AGENT_AREACODE	AGENT_PHONE	AGENT_LNAME	AGENT_YTD_SLS
-99	000	000-0000	None	\$0.00

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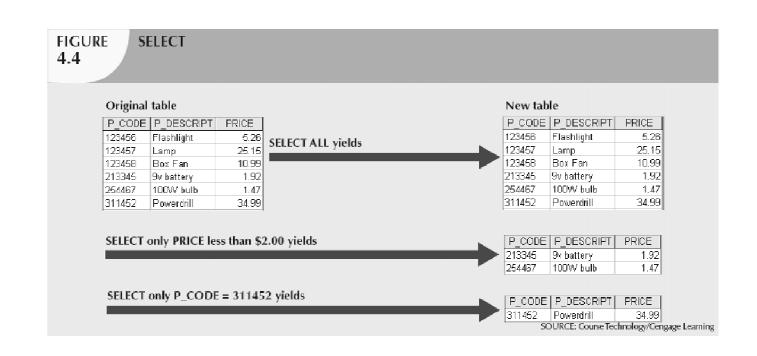
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Relational Set Operators

- Relational algebra
 - Defines theoretical way of manipulating table contents using relational operators
 - Use of relational algebra operators on existing relations produces new relations:
 - SELECT
 - PROJECT
 - JOIN
 - INTERSECT

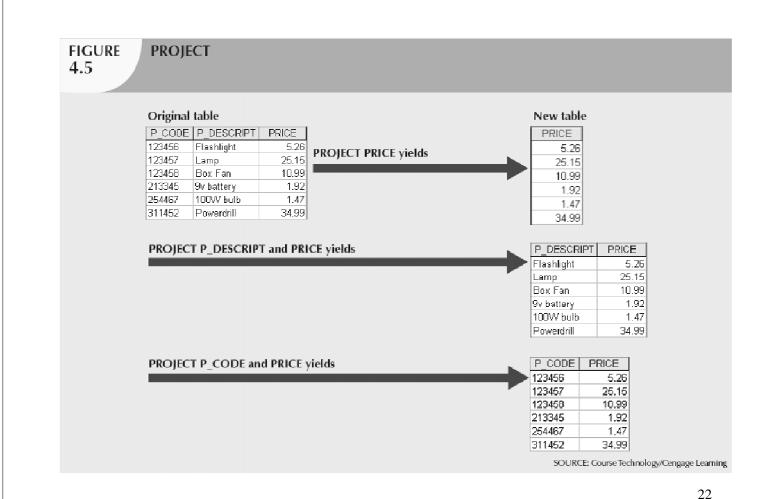
- UNION
- DIFFERENCE

- PRODUCT
- DIVIDE

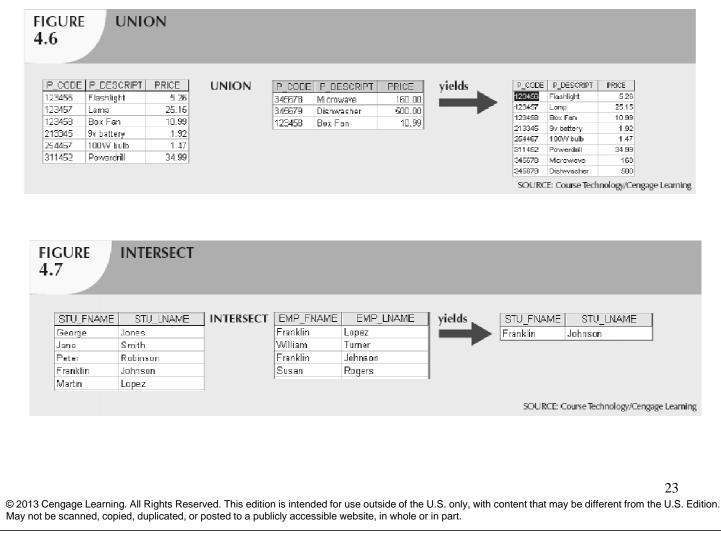


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.8	Έ							
STU_FNAME STU_LNAME George Jones Jane Smith Peter Robinson Franklin Johnson Martin Lopez	DIFFERENCE	EMP_FNAN Franklin William Franklin Susan	IE EMP. Lopez Turner Johnson Rogers	LNAME	yields	STU_FNAN George Jane Peter Martin	/E STU_LNA Jones Smith Robinson Lopez CE: Course Technol	 gage Learni
IGURE PRODUCT								
P_CODE P_DESCRIPT PRICE	11000001	STORE AISLE	SHELF	yields 🛌	P CODE		PRICE STORE	

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Relational Set Operators (cont'd.)

- Natural join
 - Links tables by selecting rows with common values in common attributes (join columns)
- Equijoin
 - Links tables on the basis of an equality condition that compares specified columns
- Theta join
 - Any other comparison operator is used

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Relational Set Operators (cont'd.)

- Inner join
 - Only returns matched records from the tables that are being joined
- Outer join
 - Matched pairs are retained, and any unmatched values in other table are left null



Two tables that will be used in join illustrations

Table name: CUSTOMER

CUS_LNAME	CUS_ZIP	AGENT_CODE
Walker	32145	231
Adares	32145	125
Rakowski	34129	167
Rodriguez	37134	125
Smithson	37134	421
Vanioo	32145	231
	Valker Adares Rakowski Rodriguez Smithson	Walker 32145 Adares 32145 Rakowski 34129 Rodriguez 37134 Smithson 37134

Table name: AGENT

AGENT_CODE	AGENT_PHONE
125	6152439887
167	6153426778
231	6152431124
333	9041234445

SOURCE: Course Technology/Cengage Learning

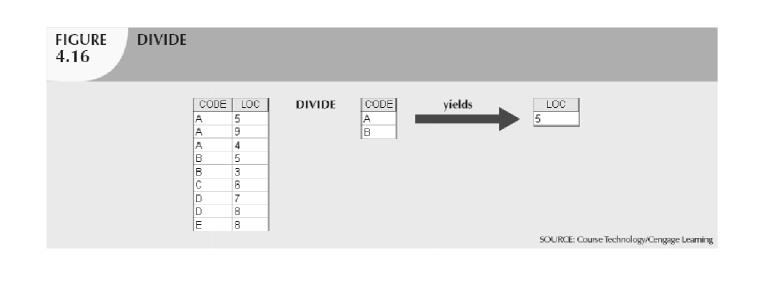
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Relational Set Operators (cont'd.)

- Left outer join
 - Yields all of the rows in the CUSTOMER table
 - Including those that do not have a matching value in the AGENT table
- Right outer join
 - Yields all of the rows in the AGENT table
 - Including those that do not have matching values in the CUSTOMER table



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The Data Dictionary and System Catalog

- Data dictionary
 - Provides detailed accounting of all tables found within the user/designer-created database
 - Contains (at least) all the attribute names and characteristics for each table in the system
 - Contains metadata: data about data
- System catalog
 - Contains metadata
 - Detailed system data dictionary that describes all objects within the database

TABLE 4.6 A S	ample Data Dictiona	ıry						
TABLE NAME	ATTRIBUTE NAME	CONTENTS	TYPE	FORMAT	RANGE	REQUIRED	PK or FK	FK REFERENCED TABLE
CUSTOMER	CUS_CODE	Customer account code	CHAR(5)	99999	10000-99999	Y	PK	
	CUS_LNAME	Customer last name	VARCHAR(20)	Χχοχοχοχοχ		Υ	FK	AGENT_CODE
	CUS_FNAME	Customer first name	VARCHAR(20)	Xxxxxxxxx		γ		
	CUS_INITIAL	Customer initial	CHAR(1)	Х				
	CUS_RENEW_DATE	Customer insurance	DATE	dd-mmm-yyyy				
		renewal date	CHAR(3)	999				
	AGENT_CODE	Agent code						
AGENT	AGENT_CODE	Agent code	CHAR(3)	999		γ	PK	
	AGENT_AREACODE	Agent area code	CHAR(3)	999		Υ		
	AGENT_PHONE	Agent telephone number	CHAR(8)	999-9999		Y		
	AGENT_LNAME	Agent last name	VARCHAR(20)	Xxxxxxxxx		γ		
	AGENT_YTD_SLS	Agent year-to-date sales	NUMBER (9,2)	9,999,999.99				

FK	=Foreign key
PK	=Primary key
CHAR	=Fixed character length data (1-255 characters)
VARCHAR	=Variable character length data (1-2,000 characters)
NUMBER	=Numeric data (NUMBER(9,2)) are used to specify numbers with two decimal places and up to nine digits, including the decimal places. Some
	RDBMSs permit the use of a MONEY or CURRENCY data type.

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The Data Dictionary and System Catalog (cont'd.)

- Homonym
 - Indicates the use of the same name to label different attributes
- Synonym
 - Opposite of a homonym
 - Indicates the use of different names to describe the same attribute

Relationships within the Relational Database

- 1:M relationship
 - Relational modeling ideal
 - Should be the norm in any relational database design
- 1:1 relationship
 - Should be rare in any relational database design

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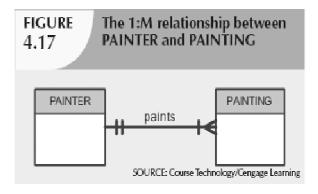
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Relationships within the Relational Database (cont'd.)

- M:N relationships
 - Cannot be implemented as such in the relational model
 - M:N relationships can be changed into 1:M relationships

The 1:M Relationship

- Relational database norm
- Found in any database environment



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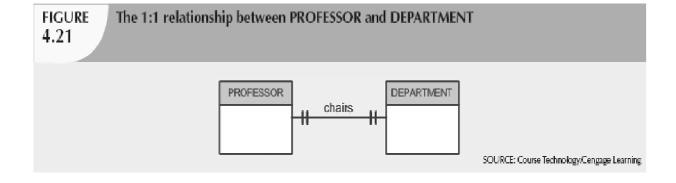
FIGURE The implemented 1:M relationship between PAINTER and PAINTING 4.18 Table name: PAINTER Primary key: PAINTER_NUM Database name: Ch04_Museum Foreign key: none PAINTER_NUM PAINTER_LNAME PAINTER_FNAME PAINTER_INITIAL 123 Ross. Georgette Ρ 126 Itero Julio G Table name: PAINTING Primary key: PAINTING NUM Foreign key: PAINTER_NUM PANTING_NUM PAINTING_TITLE PAINTER_NUM 1338 Dawn Thunder 123 1339 Vanilla Roses To Nowhere 123 1340 Tired Flounders 126 1341 Hasty Exit 123 1342 Plastic Paradise 126SOURCE: Course Technology/Cengage Learning

The 1:1 Relationship

- One entity related to only one other entity, and vice versa
- Sometimes means that entity components were not defined properly
- Could indicate that two entities actually belong in the same table
- Certain conditions absolutely require their use

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The M:N Relationship

- Implemented by breaking it up to produce a set of 1:M relationships
- Avoid problems inherent to M:N relationship by creating a composite entity
 - Includes as foreign keys the primary keys of tables to be linked

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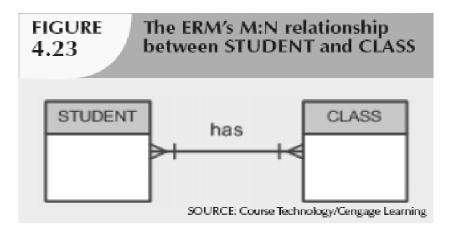


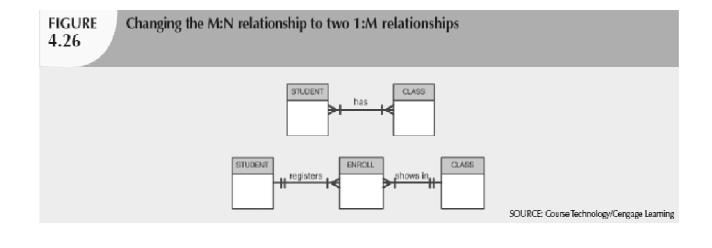
FIGURE 4.25	Converting the M:N relationship into two 1:M relationships

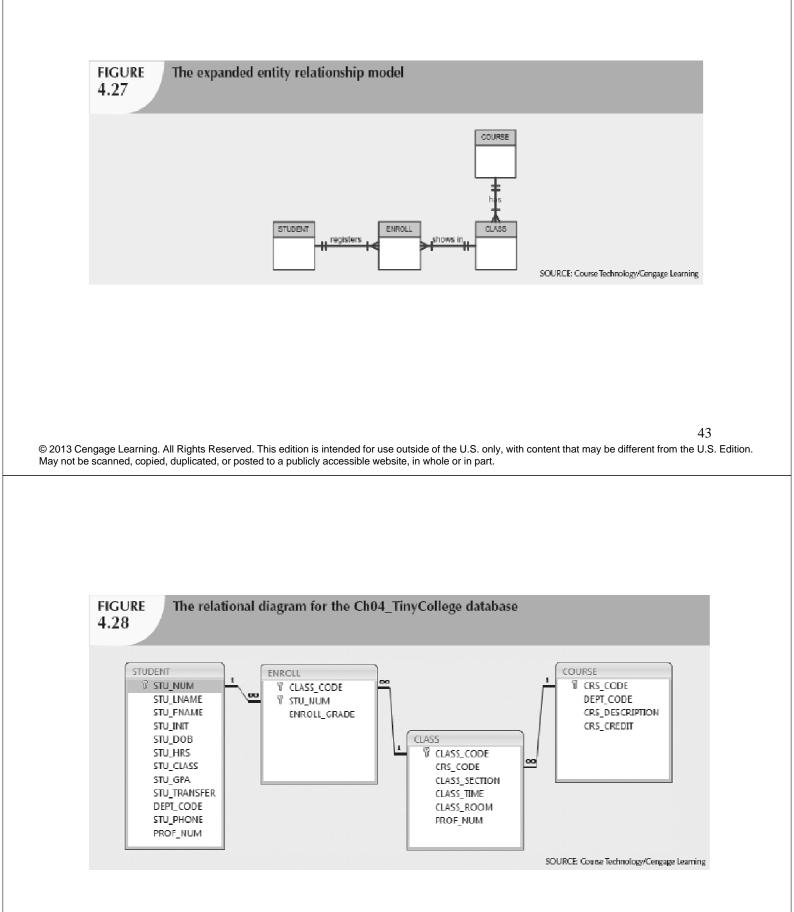
Table name: 5 Primary key: 1 Foreign key: 1 STU_NLM 5 321452 B 324257 Sr	STU_NUM none STU_LNAME owser			Databa	ase name: Ch	04_CollegeTry2			
Table name: F	ENROLL								
Primary key:	Primary key: CLASS CODE + STU NUM								
	reign key: CLASS_CODE, STU_NUM								
		ENROLL_GRADE							
10014	321452								
10014	324257	8							
10018	321452	A							
10018	324257	8							
10021	321452	C							
10021	324257	C							
Table name: (Primary key: Foreign key: (CLASS_CO	DE							
CLASS_CCDE	CRS_CODE	CLASS_SECTION	CLASS_TIME	CLASS_ROOM	PROF_NUM				
10014	ACCT-211	3	1Th 2:30-3:45 p.m.	BUS252	342				
10016	CIS-220	2	MVF 9:00-9:50 a.m.	KLR211	114				
10021	QM-261	1	MAF 8:00-8:50 a.m.	KLR200	114				
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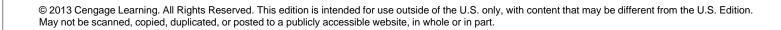
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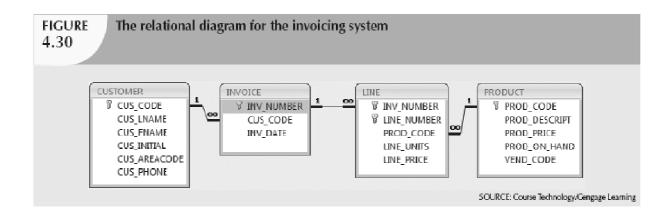




Data Redundancy Revisited

- Data redundancy leads to data anomalies
 - Can destroy the effectiveness of the database
- Foreign keys
 - Control data redundancies by using common attributes shared by tables
 - Crucial to exercising data redundancy control
- Sometimes, data redundancy is necessary



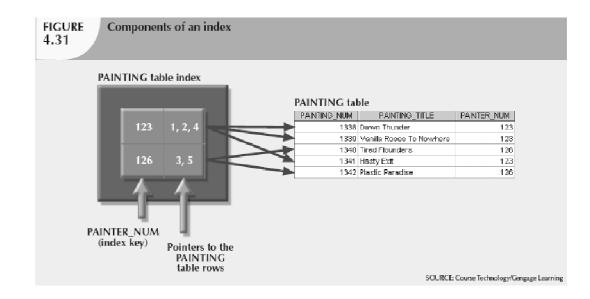


Indexes

- Orderly arrangement to logically access rows in a table
- Index key
 - Index's reference point
 - Points to data location identified by the key
- Unique index
 - Index in which the index key can have only one pointer value (row) associated with it
- Each index is associated with only one table

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Codd's Relational Database Rules

- In 1985, Codd published a list of 12 rules to define a relational database system
 - Products marketed as "relational" that did not meet minimum relational standards
- Even dominant database vendors do not fully support all 12 rules

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Summary

- Tables are basic building blocks of a relational database
- Keys are central to the use of relational tables
- Keys define functional dependencies
 - Superkey
 - Candidate key
 - Primary key
 - Secondary key
 - Foreign key

Summary (cont'd.)

- Each table row must have a primary key that uniquely identifies all attributes
- Tables are linked by common attributes
- The relational model supports relational algebra functions
 - SELECT, PROJECT, JOIN, INTERSECT UNION, DIFFERENCE, PRODUCT, DIVIDE
- Good design begins by identifying entities, attributes, and relationships

- 1:1, 1:M, M:N

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